



RAMA UNIVERSITY

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FACULTY OF ENGINEERING & TECHNOLOGY

BCS-501 Operating System

Lecturer-21

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MEMORY MANAGEMENT

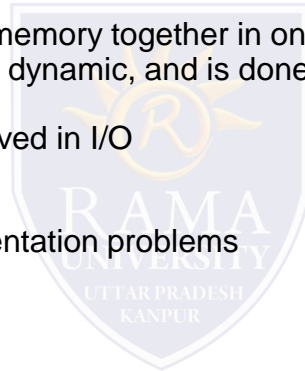
- **Fragmentation**
- **Segmentation**
- **Logical View of Segmentation**
- **Segmentation Architecture**



Fragmentation

- External Fragmentation – total memory space exists to satisfy a request, but it is not contiguous
- Internal Fragmentation – allocated memory may be slightly larger than requested memory; this size difference is memory internal to a partition, but not being used
- First fit analysis reveals that given N blocks allocated, $0.5 N$ blocks lost to fragmentation
1/3 may be unusable -> 50-percent rule
- Reduce external fragmentation by compaction
 - Shuffle memory contents to place all free memory together in one large block
 - Compaction is possible only if relocation is dynamic, and is done at execution time
 - I/O problem
 - Latch job in memory while it is involved in I/O
 - Do I/O only into OS buffers

Now consider that backing store has same fragmentation problems

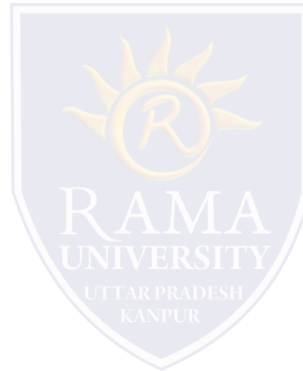


Segmentation

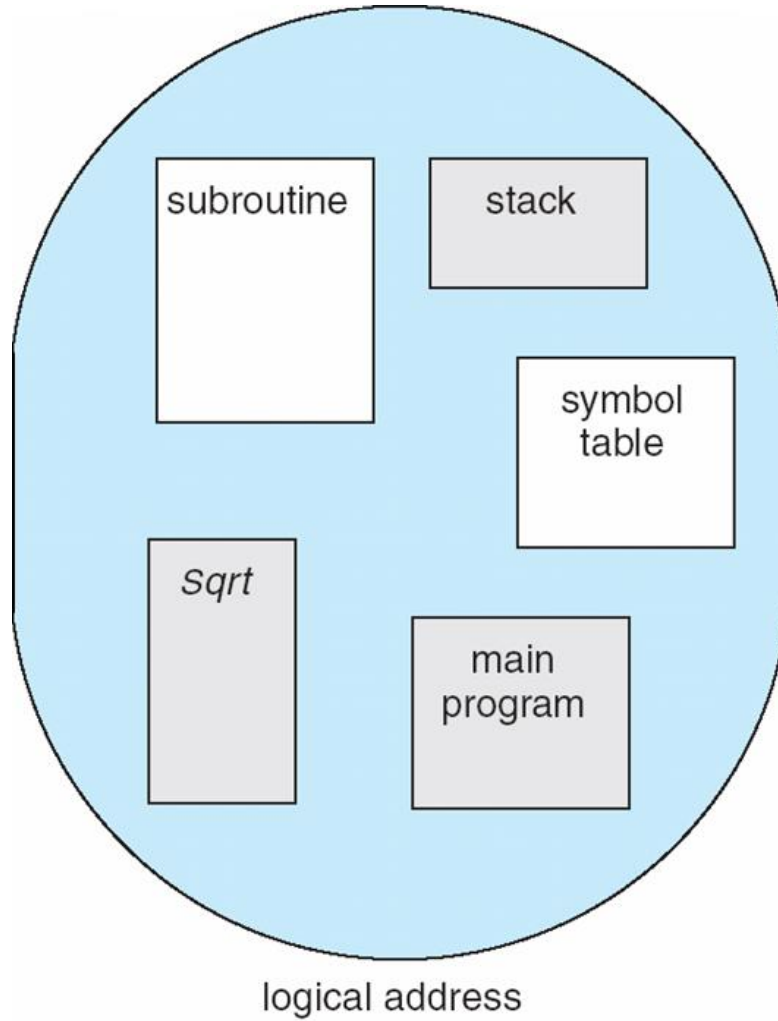
- Memory-management scheme that supports user view of memory
- A program is a collection of segments

A segment is a logical unit such as:

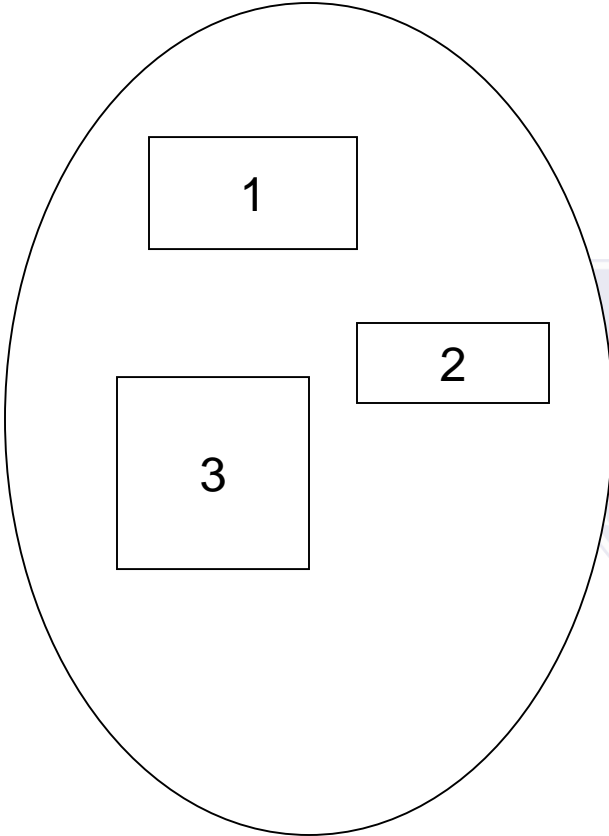
- main program
- procedure
- Function
- Method
- Object
- local variables, global variables
- common block
- Stack
- symbol table
- arrays



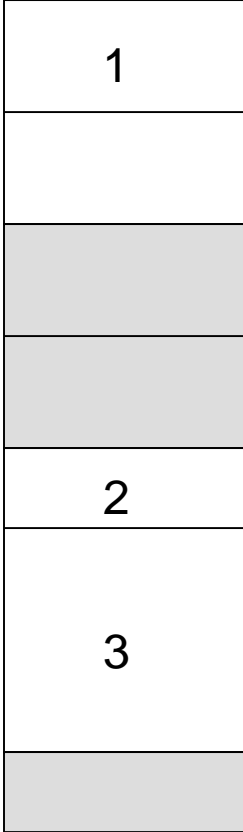
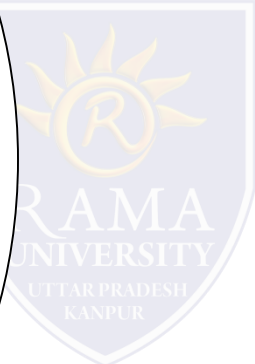
User's View of a Program



Logical View of Segmentation



user space



physical memory space

Segmentation Architecture

- Logical address consists of a two tuple:
 <segment-number, offset>
 - Segment table – maps two-dimensional physical addresses; each table entry has:
 base – contains the starting physical address where the segments reside in memory
 limit – specifies the length of the segment
 - Segment-table base register (STBR) points to the segment table's location in memory
 - Segment-table length register (STLR) indicates number of segments used by a program;
 segment number s is legal if $s < \text{STLR}$
 - Protection
 With each entry in segment table associate:
 validation bit = 0 \Rightarrow illegal segment
 read/write/execute privileges
- Protection bits associated with segments; code sharing occurs at segment level
Since segments vary in length, memory allocation is a dynamic storage-allocation problem
A segmentation example is shown in the following diagram

With relocation and limit registers, each logical address must be _____ the limit register.

- A. less than
- B. equal to
- C. greater than
- D. None of these

A process is thrashing if:

- A. it is spending more time paging than executing
- B. it is spending less time paging than executing
- C. page fault occurs
- D. swapping can not take place



When memory is divided into several fixed sized partitions, each partition may contain _____.

- A. exactly one process
- B. atleast one process
- C. multiple processes at once
- D. None of these

In fixed sized partition, the degree of multiprogramming is bounded by _____.

- A. the number of partitions
- B. the CPU utilization
- C. the memory size
- D. All of these

Transient operating system code is code that

- A. is not easily accessible
- B. comes and goes as needed
- C. stays in the memory always
- D. never enters the memory space

